# Programming (Econometrics)

Lecture 5: Linear data structures

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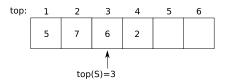
### Array

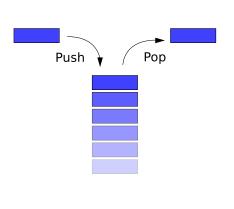
- Data structures allow to store a set of elements and guarantee certain complexity for elementary operations (access, insert, delete, search x, search min/max)
- Arrays are the most elementary data structures
- Random access: O(1)
- All other operations: O(n)



### Stack

- Last-In First-Out access semantics (LIFO)
- Can be implemented using an array and index of top element







#### Use of stacks: call stack

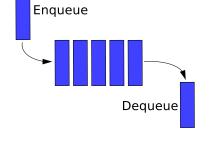
■ When you make a method call, the new frame of execution (local variables) is pushed to the stack

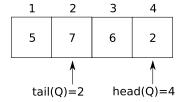
When the method exits, the local variables are simply popped from the stack



### Queue

- First-In First-Out access semantics (FIFO)
- Can be implemented using an array and indices of first (head) and last (tail) elements
- Empty queue is denoted with head = 0, tail = 1





4 queues (3, 7, 6, 2), 3 dequeues (delete 3, 7 and 6), and 1 enqueue (5): contents of Q are [2, 5].



# Use of queues: simulation with (multiple) queues

Many practical problems can be modeled with queues

■ E.g. factory arriving material to be processed, process step 1 storage, process step 2 storage, ...



### Dynamic data structures

 Until now all the data structures we considered have been static

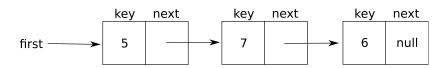
• When elements are constantly inserted / deleted, static structures are slow (O(n))

Need for node-based dynamic structures



#### Linked list

- Linked list is a list where each element has its own node, that contains the key and a reference to the next node
- Can be used to implement a stack





# Pass by reference in Matlab with OO

■ OO-extensions and their use in passing by reference

```
classdef node < handle
  properties
    key
    next
  end
end</pre>
```



#### Linked list construction

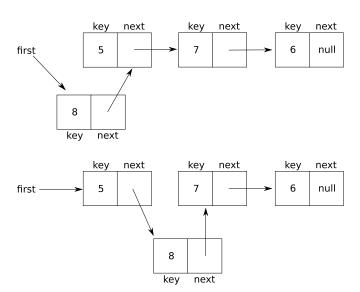
```
classdef linkedlist < handle
  properties
    first
  end
end
function L = initLinkedList(value)
  L = linkedList();
  fNode = node();
  fNode.key = value;
  fNode.next = [];
  L. first = fNode;
end
```



#### Linked list traversal

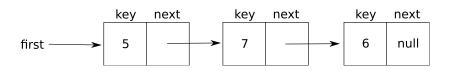
```
function node = findKey(L, key)
  curNode = L.first;
  while (~isempty(curNode))
    if (curNode.key == key)
      node = curNode;
      break;
    end
    curNode = curNode.next:
  end
end
            key
                 next
                            key
                                 next
                                            key
                                                 next
                             7
                                            6
             5
                                                 null
first
```

#### Linked list insert element





#### Linked list delete element

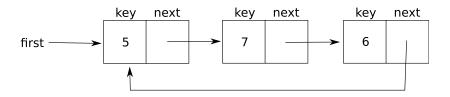


#### deleteNodeAfter(nodeOf5)





## Circular linked list



■ Uses: round-robin scheduling of processes in multi-tasking environments (e.g. your computer)



# Complexity of linked list operations

- Insert/delete element in beginning: O(1)
- Insert/delete element at current iteration location: O(1)
- Random access: O(n)
- Search: O(n)

prev ->next = toDelete ->next;
delete toDelete;
// if only forgetting were









